

# Lecture 1: Foundations of ontology engineering

home: <http://www.meteck.org>

*Foundations and recent trends on ontology engineering*  
*Universitat Politècnica de Catalunya, 2017*

# Outline

- 1 Introduction
- 2 Where is it used?
- 3 What is an Ontology?
- 4 Logic and automated reasoning
  - Representation languages
  - Automated reasoning

# Administrivia (1/2)

- This course consists of two lectures of four hours, exercises, and a mini-project assignment in small groups
- Labs&self study about 4-6 hours
- Mini-project [100%], 8 hours (rough timeframe)
  - Topics will be distributed after this lecture
  - Need to form groups of 2-3 people and choose topic by next lecture
  - Deadline hand-in: 14 June
- Following the lectures will be easier and beneficial when you have read the recommended and required reading beforehand, and at least the 10-20 pages/chapter of the lecture notes

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  - The slides serve as a teaching aid, not as a neat summary

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- ... so there is no single textbook (yet) that covers all topics for the novice ontologist, has exercises with given, clear answers etc, but...

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- ... there are lecture notes (though you still have to read some scientific literature)

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# An ontology (very informally)

- classes, relationships between them, and constraints that hold between/for them, with possibly individuals and their relations
- as a representation of a particular subject domain



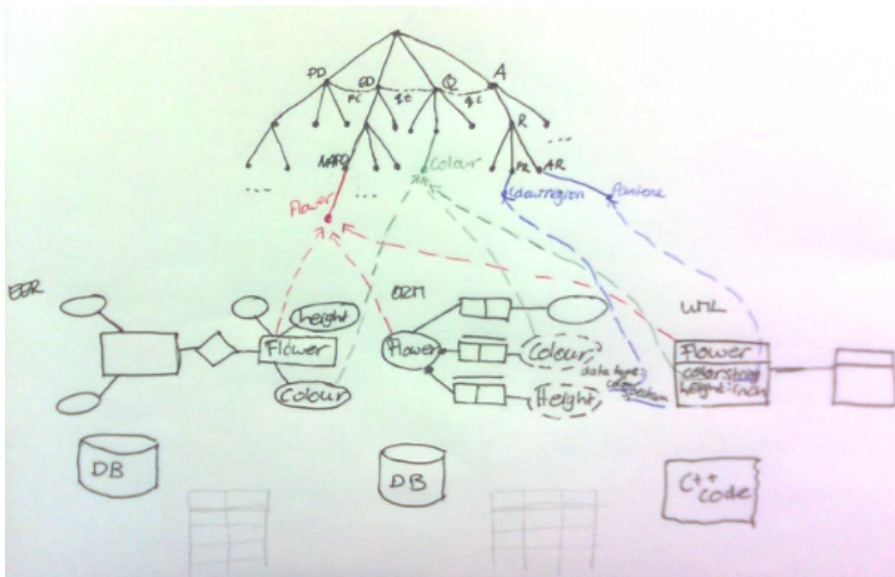


# Conceptual data models vs ontologies

- Main differences:
  - Information needs for one application vs. **representing the knowledge of a subject domain** (regardless the particular application)
  - **Formalization** in a logic language (though one could do that for conceptual models as well)

# Conceptual data models vs ontologies

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  - **Formalization** in a logic language (though one could do that for conceptual models as well)
- An ontology as a layer on top of conceptual data models
  - To improve the quality of a conceptual data model (hence, the software)
  - To facilitate database integration, or prevent the usual data integration problems



# Databases vs. Knowledge bases

- Main differences:
  - Representation of the knowledge
  - Rules
  - Reasoning to infer new or implicit knowledge, detect inconsistencies of the knowledge base
  - Open World Assumption (vs. Closed World Assumption in databases)

# What is the usefulness of an ontology?

- Making, more or less precisely, the (dis-)agreement among people explicit
- Enrich software applications with the additional semantics  $\Rightarrow$  *ontology-driven information systems*
- Thus, practically, improving computer-computer, computer-human, and human-human communication

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# Examples ontologies in information systems

- **e-learning** with *Inquire Biology* [Chaudhri et al., 2013]: textbook annotated with terms of the ontology, generates questions and answers.
- **data integration, cultural heritage**: combining resources of data and querying them, on the ever interesting topic of food [Calvanese et al., 2016]
- **publishing** of scientific papers, books: enable navigation and understanding of scholarly documents [Di lorio et al., 2014]
- **semantic meta-mining of data mining experiments** (sections 1 and 5 of [Keet et al., 2015]): mine the (ontology-based) annotations of the data mining experiments, reason over that to have it propose the optimal data mining experiment



# More Examples

- **For science** Inside the scientific method: Outperforming humans (ontology+reasoner): classification of protein phosphatases [Wolstencroft et al., 2007]
- **Deep Question-Answering** with Watson beating human top-performers in 'Jeopardy!'; uses over 100 techniques, including ontologies for integration
- **Ontology-driven conceptual data modelling**: being more precise than just drawing diagrams, e.g., on those 'shared' and 'composite' aggregations in UML Class diagrams [Keet & Artale, 2008], finding contradictions.

# The Semantic Web – Introduction

## (some motivations for ontologies and knowledge bases)

- AI put to the test in the (uncontrollable?) very large field
- Adding meaning to plain HTML pages and Web 2.0 by using theory and technologies of KBs and ontologies

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### Homepage of Maria Keet

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Research
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CV, short bio, **publications**, and my Google Scholar profile

*\*\*\*If you ended up on this page, but are not interested in my work, then choose **Other topics**, where there is information on Escher, book reviews, some photos and other not-quite-random themes.\*\*\**



# The Semantic Web – Introduction

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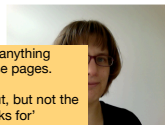
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### Homepage of Maria Keet



#### Dr. C. Maria Keet

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Blog: keet blog  
email: mkeet at cs.uct.ac.za  
linkedIn



The plain Web's "a href" doesn't say anything about how the webpage links to those pages.

Implicitly (what humans can figure out, but not the computer), we have 'Academic' 'works for' 'University' (and, more generally: employee works for organisation).

-> add such information, find things more easily

```
<b>Dr. C. Maria Keet</b><br>
Senior Lecturer<br>
<a href="http://www.cs.uct.ac.za/">Department of Computer Science</a><br>
<a href="http://www.uct.ac.za/">University of Cape Town</a><br>
```

pose Other topics, where te-random themes.\*\*\*

# Generalising from the examples, ontologies are used for:

- Data(base) linking and integration
- Instance classification
- Matchmaking and services
- Querying, information retrieval
  - Ontology-Based Data Access
  - Ontologies to improve NLP
- Bringing more quality criteria into conceptual data modelling to develop a better model (hence, a better quality software system)

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# Background

- Aristotle and colleagues: **Ontology**
- Engineering: ontologies (count noun)
- Investigating reality, representing it
- Putting an engineering artefact to use

What then, is this engineering artefact?



(Guarino, 2002)



First, let's look at an artefact: a text file....



```
<owl:Class rdf:about="&AfricanWildlifeOntology1;lion">
  <rdfs:subClassOf rdf:resource="&AfricanWildlifeOntology1;animal"/>
  <rdfs:subClassOf>
    <owl:Restriction>
      <owl:onProperty rdf:resource="&AfricanWildlifeOntology1;eats"/>
      <owl:allValuesFrom rdf:resource="&AfricanWildlifeOntology1;herbivore"/>
    </owl:Restriction>
  </rdfs:subClassOf>
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    <owl:Restriction>
      <owl:onProperty rdf:resource="&AfricanWildlifeOntology1;eats"/>
      <owl:someValuesFrom rdf:resource="#Impala"/>
    </owl:Restriction>
  </rdfs:subClassOf>
  <rdfs:comment>Lions are animals that eat only herbivores.</rdfs:comment>
</owl:Class>

<!-- file://Applications/Protege_4.1_beta/AfricanWildlifeOntology1.owl#plant -->

<owl:Class rdf:about="&AfricanWildlifeOntology1;plant">
  <rdfs:comment>Plants are disjoint from animals.</rdfs:comment>
</owl:Class>
```

... and rendered in an ontology editor

The screenshot shows the AfricanWildlifeOntology2 editor interface. The browser address bar displays the URL: <http://www.meteck.org/teaching/ontologies/AfricanWildlifeOntology2.owl>. The main menu includes: Active Ontology, Entities, **Classes**, Object Properties, Data Properties, Individuals, OWLViz, DL Query, and OntoGraf.

The **Classes** tab is active, showing the **Class hierarchy: lion**. The hierarchy is as follows:

- animal
  - Impala
  - Omnivore
  - RockDassie
  - Warthog
  - carnivore
  - giraffe
  - herbivore
  - lion** (selected)
- plant
  - CarnivorousPlant
  - Grass
  - Palmtree
  - tasty-plant
  - tree

The **Annotations: lion** panel shows the following annotations:

- comment**: "Lions are animals that eat only herbivores."

The **Description: lion** panel shows the following descriptions:

- Equivalent classes**: (empty)
- Superclasses**:
  - animal
  - eats **only** herbivore
  - eats **some** Impala

At the bottom, there is a status bar with the text: "To use the reasoner click Reasoner->Start reasoner" and a checked checkbox for "Show Inferences".

# A few definitions on what the text in the file is supposed to stand for

- Most cited (but very inadequate definition): “An ontology is a specification of a conceptualization” (by Tom Gruber, 1993)
- “a formal specification of a shared conceptualization” (by Borst, 1997)
- “An ontology is a formal, explicit specification of a shared conceptualization” (Studer et al., 1998)

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- “a formal specification of a shared conceptualization” (by Borst, 1997)
- “An ontology is a formal, explicit specification of a shared conceptualization” (Studer et al., 1998)
- What is a *conceptualization*, and a *formal, explicit specification*? Why *shared*?

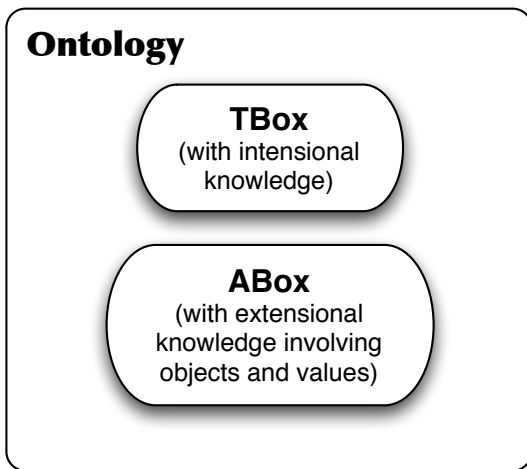
## More definitions

- More detailed: “An ontology is a logical theory accounting for the *intended meaning* of a formal vocabulary, i.e. its *ontological commitment* to a particular *conceptualization* of the world. The intended models of a logical language using such a vocabulary are constrained by its ontological commitment. An ontology indirectly reflects this commitment (and the underlying conceptualization) by approximating these intended models.” (Guarino, 1998)

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- And back to a simpler definition: “with an ontology being equivalent to a [Description Logic knowledge base](#)” (Horrocks et al, 2003)

## Description Logic knowledge base



# From logical to ontological level (1/2)

- Logical level (no structure, no constrained meaning<sup>1</sup>):
  - $\exists x(Apple(x) \wedge Green(x))$
  - “there exists an object that is an apple and it is green”

*adapted from (Guarino, 2008)*

<sup>1</sup> meaning in the sense of subject domain semantics, not formal semantics

<sup>2</sup> DL has a model-theoretic semantics, so the axioms have a meaning in that sense of ‘meaning/semantics’



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  - $\exists x : apple\ Green(x)$  (many-sorted logics)
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  - $Apple(a)$  and  $hasColor(a, green)$  (description logics<sup>2</sup>)
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  - $\cancel{Green(a)}$  and  $\cancel{hasShape(a, apple)}$
  - “object  $a$  is a green and that object  $a$  has the shape of an apple”

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- Ontological level (structure, constrained meaning):
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  - e.g., 'apple objects' seems better than 'green objects'
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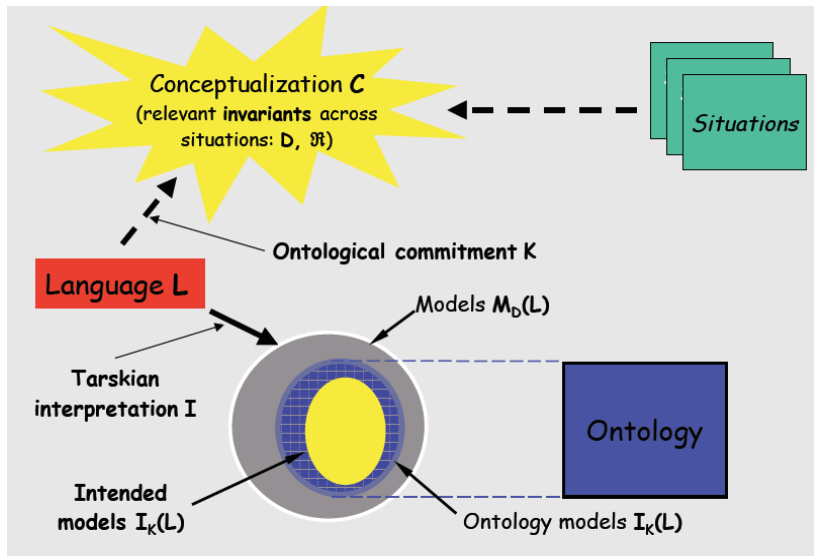
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- Put differently: one way of representing things turn out to be *better* than others.

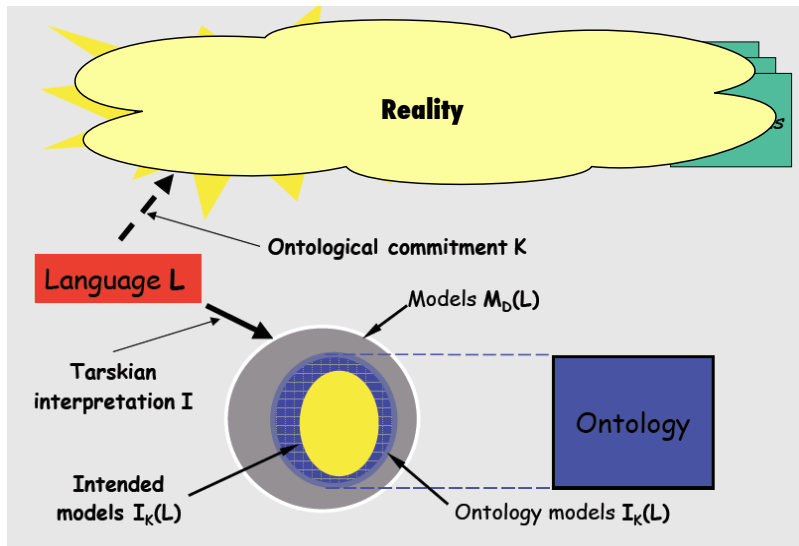
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# Ontologies and meaning

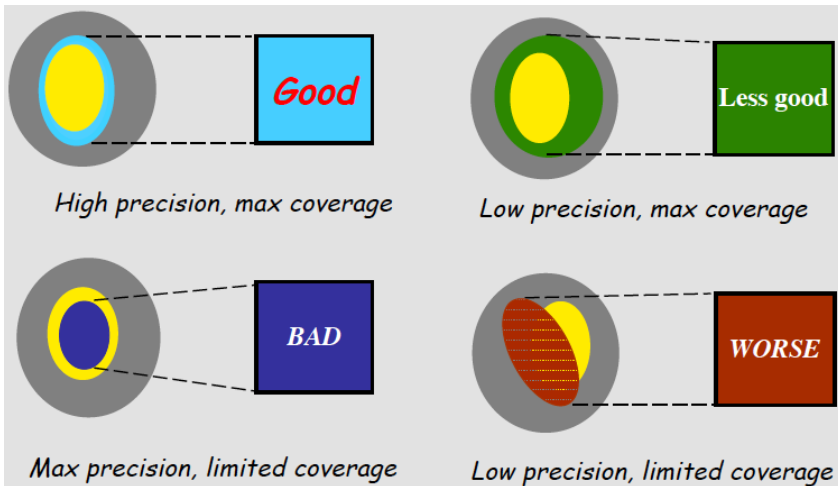




# Ontologies and reality



# Quality of the ontology



(Guarino, 2002)

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# Preliminary note

- There are only a few core concepts to get the general idea
- There are very many details
- Here we focus on the core concepts and some details and how that works out in computing
- More logic and details in the lecture notes

# How to formalise it?

- Logics have a:
  - Syntax
    - Alphabet
    - Languages constructs
    - Sentences to assert knowledge
  - Semantics
    - Formal meaning

## Several ontology languages

- W3C-standardised Web Ontology Language OWL, comes in many 'species'
  - Description Logics-based OWL species
  - OWL full and OWL 2 full (RDF-based semantics)
- Common logic, CLIF
- First order logic
- Fuzzy and temporal extensions and variants

## DLs are structured fragments of FOL

- Recall that full FOL is undecidable
- This is unpleasant for automated reasoning





# DLs are structured fragments of FOL

- We end up with *trade-offs* of features in a DL
- Some features always will make the language undecidable (e.g., true role composition,  $R \circ S \equiv T$ )
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- E.g., one could define a language where:
  - it is prohibited to use  $\neg$  (negation) in an axiom, or
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  - $\exists R$  only on the rhs of the inclusion but not on the lhs
- There are *many* DLs, and most combinations have been investigated over the past 25 years
- Roughly: the fewer features and the more restrictions, the more ‘computationally well-behaved’ the language is

# On those OWL 'species'

- OWL standardised in 2004
- OWL 2 standardised in 2009

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    - OWL 2 EL
    - OWL 2 QL
    - OWL 2 RL
  - OWL 2 DL
  - OWL 2 full

# OWL 2 DL Overview

- Has more features than OWL DL
- Computationally more 'costly' (N2EXPTIME-complete cf EXPTIME-complete)
- Based on the DL language *SR<sub>Q</sub>IQ*
- Main novelty especially w.r.t. modelling practices: qualified cardinality constraints, more 'characteristics' of object properties

# OWL 2 EL Overview

- Intended for large 'simple' ontologies
- Focussed on type-level knowledge (TBox)
- Better computational behaviour than OWL 2 DL (polynomial vs. exponential/open)
- Based on the DL language  $\mathcal{EL}^{++}$  (PTime complete)
- Reasoner: e.g. CEL <http://code.google.com/p/cel/>
- $\sqcap \exists T \perp \sqsubseteq \sqcap \exists T \perp$



# OWL 2 QL Overview

- Query answering over a large amount of instances with same kind of performance as relational databases (Ontology-Based Data Access)
- Expressive features cover several used features of UML Class diagrams and ER models ('CONceptual MOdel-based Data Access')
- Based on  $DL-Lite_{\mathcal{R}}$  (*more is possible with UNA and in some implementations*)

# OWL 2 RL Overview

- Development motivated by: what fraction of OWL 2 DL can be expressed by rules (with equality)?
- Scalable reasoning in the context of RDF(S) application
- Rule-based technologies (forward chaining rule system, over *instances*)
- Inspired by Description Logic Programs and pD\*
- Reasoning in PTime
- No  $\forall$  and  $\neg$  on lhs, and  $\exists$  and  $\sqcup$  on rhs of  $\sqsubseteq$

# OWL Syntax—Many notations, actually

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  - e.g., not “ $\sqsubseteq$ ” to process inclusion, but “SubClassOf”

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- Making those DL symbols usable by a computer
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- RDF/XML
  - Official exchange syntax
  - Hard for humans to read (and RDF parsers are hard to write)
- OWL/XML
  - Not the RDF syntax
  - Still hard for humans, but more XML than RDF tools available
- Abstract syntax
  - To some, considered human readable
- “User-usable” ones
  - e.g., Manchester syntax, informal and limited matching with UML, pseudo-NL verbalisations

# Example correspondences

- 'Each C is a D' / 'All Cs are Ds'
    - `SubClassOf(C D)`
    - $C \sqsubseteq D$
    - $\forall x(C(x) \rightarrow D(x))$
- e.g.: *Giraffe*  $\sqsubseteq$  *Animal*

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- 'Each C R at least one D' / 'Each C R some D'
  - `SubClassOf(C ObjectSomeValueFrom(R D))`
  - $C \sqsubseteq \exists R.D$
  - $\forall x(C(x) \rightarrow \exists y(R(x,y) \wedge D(y))$
  - e.g.: *Elephant*  $\sqsubseteq \exists \text{eats.AmarulaFruit}$

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  - e.g.: *Elephant*  $\sqsubseteq \exists \text{eats.AmarulaFruit}$
- 'C and D are disjoint' / 'each C is not a D'
  - `DisjointClasses(C D) /`  
`SubClassOf(C ObjectComplementOf(D))`
  - $C \sqcap D \sqsubseteq \perp$  (disj.) /  $C \sqsubseteq \neg D$  (complement)
  - $\forall x(C(x) \rightarrow \neg D(x))$
  - e.g.: *Herbivore*  $\sqsubseteq \neg \text{Carnivore}$

# Semantics (DL-based OWL species)

- Model-theoretic semantics



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- Model-theoretic semantics
- Domain  $\Delta$  is a non-empty set of objects
- Interpretation:  $\cdot^{\mathcal{I}}$  is the *interpretation function*, domain  $\Delta^{\mathcal{I}}$ 
  - $\cdot^{\mathcal{I}}$  maps every concept name  $A$  to a subset  $A^{\mathcal{I}} \subseteq \Delta^{\mathcal{I}}$
  - $\cdot^{\mathcal{I}}$  maps every role name  $R$  to a subset  $R^{\mathcal{I}} \subseteq \Delta^{\mathcal{I}} \times \Delta^{\mathcal{I}}$
  - $\cdot^{\mathcal{I}}$  maps every individual name  $a$  to elements of  $\Delta^{\mathcal{I}}$ :  $a^{\mathcal{I}} \in \Delta^{\mathcal{I}}$
- Note:  $\top^{\mathcal{I}} = \Delta^{\mathcal{I}}$  and  $\perp^{\mathcal{I}} = \emptyset$

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- Interpretation:  $\cdot^{\mathcal{I}}$  is the *interpretation function*, domain  $\Delta^{\mathcal{I}}$ 
  - $\cdot^{\mathcal{I}}$  maps every concept name  $A$  to a subset  $A^{\mathcal{I}} \subseteq \Delta^{\mathcal{I}}$
  - $\cdot^{\mathcal{I}}$  maps every role name  $R$  to a subset  $R^{\mathcal{I}} \subseteq \Delta^{\mathcal{I}} \times \Delta^{\mathcal{I}}$
  - $\cdot^{\mathcal{I}}$  maps every individual name  $a$  to elements of  $\Delta^{\mathcal{I}}$ :  $a^{\mathcal{I}} \in \Delta^{\mathcal{I}}$
- Note:  $\top^{\mathcal{I}} = \Delta^{\mathcal{I}}$  and  $\perp^{\mathcal{I}} = \emptyset$
- An interpretation  $\mathcal{I} = (\Delta^{\mathcal{I}}, \cdot^{\mathcal{I}})$  is a **model** of a knowledge base  $\mathcal{KB}$  if every axiom of  $\mathcal{KB}$  is satisfied by  $\mathcal{I}$
- A knowledge base  $\mathcal{KB}$  is said to be **satisfiable** if it admits a model

# Essential to automated reasoning (Ch2 of LN)

- The choice of the class of problems the software program has to solve
- The formal language in which to represent the problems
- The way how the program has to compute the solution
- How to do this efficiently

# Reasoning services for DL-based OWL ontologies

- Concept (and role) satisfiability ( $\mathcal{KB} \not\models C \sqsubseteq \perp$ )
  - is there a model of  $\mathcal{KB}$  in which  $C$  (resp.  $R$ ) has a nonempty extension?

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- Instance retrieval ( $\{a \mid \mathcal{KB} \models C(a)\}$ )
  - find all members of  $C$  in  $\mathcal{KB}$ , i.e., compute all individuals  $a$  s.t.  $C(a)$  is satisfied by every interpretation of  $\mathcal{KB}$



# Logical implication

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 $\text{Undergrad}(\text{John})$
- $\mathcal{KB} \models \text{Professor}(\text{John})?$  or perhaps  
 $\mathcal{KB} \models \neg \text{Professor}(\text{John})?$

# Automated reasoning techniques

- How do we compute, say, satisfiability?

# Automated reasoning techniques

- How do we compute, say, satisfiability?
- Truth tables are too cumbersome
- Several techniques are more efficient
- Current 'winner' is *tableau reasoning*

# The idea, same as for FOL

- A sound and complete procedure deciding satisfiability is all we need, and the **tableaux method is a decision procedure which checks the existence of a model**
- It exhaustively looks at all the possibilities, so that it can eventually prove that no model could be found for unsatisfiable formulas.
- $\phi \models \psi$  iff  $\phi \wedge \neg\psi$  is NOT satisfiable—if it is satisfiable, we have found a counterexample
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- $\phi \models \psi$  iff  $\phi \wedge \neg\psi$  is NOT satisfiable—if it is satisfiable, we have found a counterexample
- Decompose the formula in top-down fashion
- Following slide *simplified* process (thanks to Markus Krötzsch & Sebastian Rudolph ESSLLI 2009 Bordeaux)



```

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## Knowledge Base

## Tableau

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ex:HappyCatOwner


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**ex:HappyCatOwner**  
 [owl:intersectionOf (■, ■)]



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## Knowledge Base



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## Knowledge Base



ex:owns



ex:HappyCatOwner  
[owl:intersectionOf ([■], [■])]  
■ ■

ex:Cat

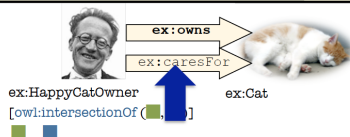


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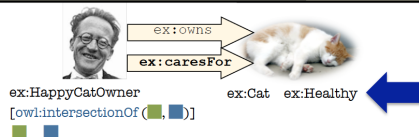
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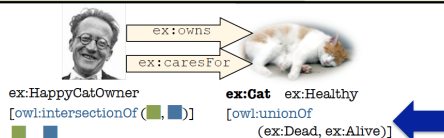
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## Knowledge Base



ex:owns

ex:caresFor



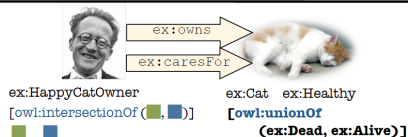
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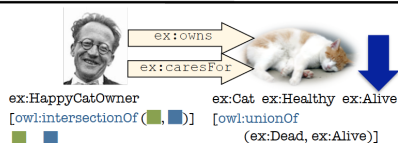
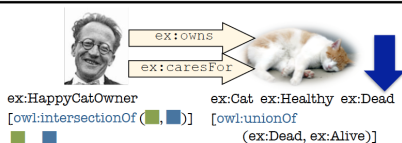
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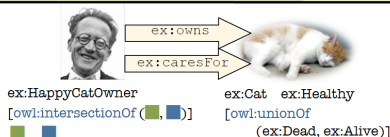
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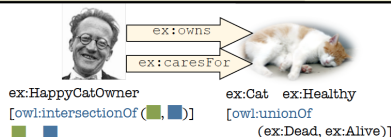


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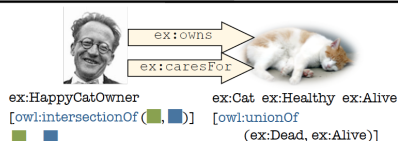
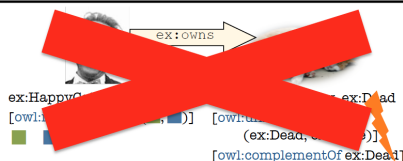


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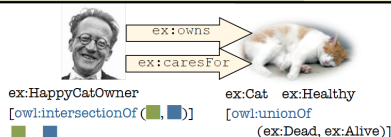


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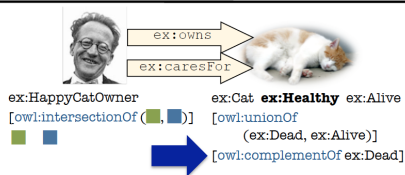
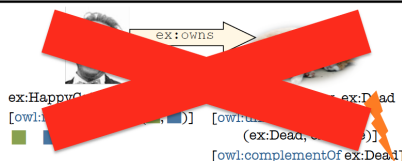
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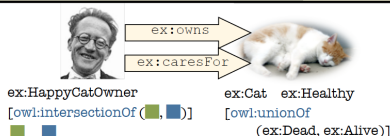
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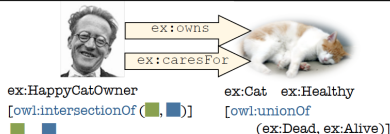
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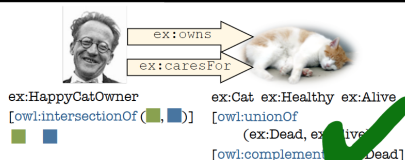
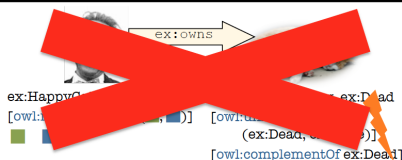
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## Knowledge Base

# Satisfiable



## Tableau



# Summary

- 1 Introduction
- 2 Where is it used?
- 3 What is an Ontology?
- 4 Logic and automated reasoning
  - Representation languages
  - Automated reasoning



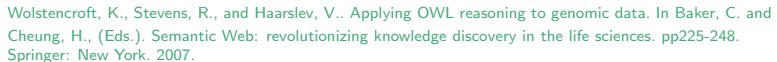
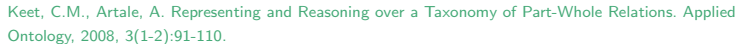
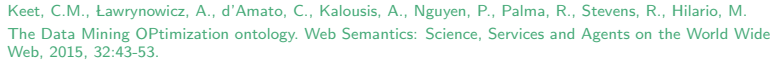
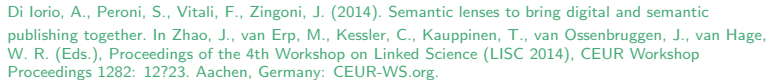
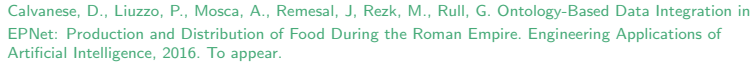
# Exercises

- Chapter 1: exercises 1 and 2
- Chapter 3: exercises 5 (optionally 6)
- Chapter 4: exercises 8, 10, 14-21
- Read the mini-project topics (optionally also looking up what some of those terms mean) and select one

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- The tutorial ontologies are available from  
<http://www.meteck.org/teaching/ontologies/>

\_\_\_\_\_



# Supported class restrictions OWL 2 EL

- existential quantification to a class expression or a data range
- existential quantification to an individual or a literal
- self-restriction
- enumerations involving a single individual or a single literal
- intersection of classes and data ranges

# Supported axioms, restricted to allowed set of class expressions OWL 2 EL

- class inclusion, equivalence, disjointness
- object property inclusion and data property inclusion
- property equivalence
- transitive object properties
- reflexive object properties
- domain and range restrictions
- assertions
- functional data properties
- keys
- In short:  $\sqcap \exists \top \perp \sqsubseteq \sqcap \exists \top \perp$

# NOT supported in OWL 2 EL

- universal quantification to a class expression or a data range
- cardinality restrictions
- disjunction
- class negation
- enumerations involving more than one individual
- disjoint properties
- irreflexive, symmetric, and asymmetric object properties
- inverse object properties, functional and inverse-functional object properties

# Supported Axioms in OWL 2 QL, restrictions

- Subclass expressions restrictions:
  - a class
  - existential quantification (ObjectSomeValuesFrom) where the class is limited to owl:Thing
  - existential quantification to a data range (DataSomeValuesFrom)
- Super expressions restrictions:
  - a class
  - intersection (ObjectIntersectionOf)
  - negation (ObjectComplementOf)
  - existential quantification to a class (ObjectSomeValuesFrom)
  - existential quantification to a data range (DataSomeValuesFrom)

## Supported Axioms in OWL 2QL

- Restrictions on class expressions, object and data properties occurring in functionality assertions cannot be specialized
- subclass axioms
- class expression equivalence (involving `subClassExpression`), disjointness
- inverse object properties
- property inclusion (not involving property chains and `SubDataPropertyOf`)
- property equivalence
- property domain and range
- disjoint properties
- symmetric, reflexive, irreflexive, asymmetric properties
- assertions other than individual equality assertions and negative property assertions (`DifferentIndividuals`, `ClassAssertion`, `ObjectPropertyAssertion`, and `DataPropertyAssertion`)



# NOT supported in OWL 2 QL

- existential quantification to a class expression or a data range in the subclass position
- self-restriction
- existential quantification to an individual or a literal
- enumeration of individuals and literals
- universal quantification to a class expression or a data range
- cardinality restrictions
- disjunction
- property inclusions involving property chains
- functional and inverse-functional properties
- transitive properties
- keys
- individual equality assertions and negative property assertions

## Supported in OWL 2 RL

- More restrictions on class expressions (see table 2, e.g. no SomeValuesFrom on the right-hand side of a subclass axiom)
- All axioms in OWL 2 RL are constrained in a way that is compliant with the restrictions in Table 2.
- Thus, OWL 2 RL supports all axioms of OWL 2 apart from disjoint unions of classes and reflexive object property axioms.
- No  $\forall$  and  $\neg$  on lhs, and  $\exists$  and  $\sqcup$  on rhs of  $\sqsubseteq$